

### 1.3 Combinatorics

Def: A hypergraph  $H = (V, E)$ , where  $V$  is a finite set (the vertices), and  $E$  is a family of subsets of  $V$  (the edges). It is  $k$ -uniform if every edge contains exactly  $k$  vertices.

Def:  $H$  has property  $\mathcal{B}$ , or it's 2-colorable, if there is a 2-coloring of  $V$  such that no edge is monochromatic.

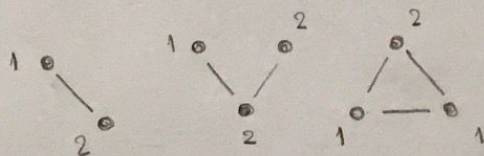
For  $k=2$  (i.e., graphs) 2-colorable is a very stringent property: There are 4 ways of coloring an edge with 2 colors, and only 2 of them satisfy  $\mathcal{B}$ . For large  $k$ , it's much easier to do it: Only 2 out of  $2^k$  colorings do not satisfy  $\mathcal{B}$ .

Def:  $m(k) = \min \{ m \mid \exists k\text{-uniform hypergraph on } m \text{ edges that doesn't have prop. } \mathcal{B} \}$

$m(1) = 1$



$m(2) = 3$



Proposition [Erdős '63]: Every  $k$ -uniform hypergraph with  $< 2^{k-1}$  edges has property  $\mathcal{B}$ . i.e.  $m(k) \geq 2^{k-1}$ .

Proof: Let  $H = (V, E)$  be a  $k$ -uniform hypergraph on  $m < 2^{k-1}$  edges.

Let  $\varphi$  be a random coloring of  $V$  with 2 colors. WTS

$\Pr[H \text{ has no monochromatic edge under } \varphi] > 0$

$\Rightarrow \Pr[H \text{ has monochromatic...}] < 1$

$\Pr[H \text{ has monochromatic...}] = \Pr[\exists e \in E : e \text{ monochromatic}]$

$$\leq \sum_{e \in E} \Pr[e \text{ monochromatic}]$$

$$= m \cdot 2 \cdot 2^{-k}$$

$$= m / 2^{k-1}$$

$$< 1 \quad (\text{by assumption}). \quad \square$$

Cherukashin & Kozim 2015:  $m(k) \geq \Omega\left(\left(\binom{k}{m} \cdot 2^k\right)^{1/2}\right)$ .

Theorem [Erdős '64]:  $m(k) \leq O(k^2 2^k)$

Proof: WTS  $\exists k$ -uniform hypergraph with  $O(k^2 2^k)$  edges doesn't have property  $\mathcal{B}$ .

Fix  $V$  with  $n$  vertices,  $n \gg k$ . Fix  $m \gg k$ . Let  $E = \{S_1, \dots, S_m\}$  be a random set of hyperedges with  $|S_i| = k$ . Let  $H = (V, E)$ . WTS

$\Pr[H \text{ doesn't have property } \mathcal{B}] > 0$

$\Rightarrow \Pr[H \text{ has property } \mathcal{B}] < 1$ .

$$\Pr(H \text{ has property } \mathcal{B}) = \Pr(\exists \varphi \text{ 2-coloring of } V \text{ such that no } e \in E \text{ is monochromatic})$$

(\*) : If  $m = \frac{n}{p} \ln 2$ , this is  $< 1$  and we're done.

$$\begin{aligned} &\leq \sum_{\varphi} \Pr(\text{no } e \text{ is monochromatic under } \varphi) \\ &= \sum_{\varphi} \prod_{i=1}^m \Pr(S_i \text{ is not monochromatic under } \varphi) \\ &\quad \text{(all edges are independent)} \\ &= \sum_{\varphi} \prod_{i=1}^m (1 - \underbrace{\Pr(S_i \text{ is monoch. under } \varphi)}_p) \\ &= 2^n \cdot (1-p)^m < 2^n e^{-pm} \quad (*) \end{aligned}$$

Fix an arbitrary coloring  $\varphi$ . Let  $a$  be the number of red vertices and  $b$  the number of blue vertices. Then

$$p = \frac{\binom{a}{k}}{\binom{n}{k}} + \frac{\binom{b}{k}}{\binom{n}{k}} = \frac{\binom{a}{k} + \binom{b}{k}}{\binom{n}{k}}, \quad a+b=n$$

For what  $a, b$  is  $p$  minimized?

Claim:  $f(x) = \binom{x}{k}$  is convex for  $x \geq k$ .

Proof:  $f(x)$  is just a polynomial of degree  $k$  with roots  $0, 1, \dots, k-1$  and main coefficient  $> 0$ .

$$f(x) = \frac{x!}{k!(x-k)!} = \frac{1}{k!} (x-(k-1))(x-(k-2)) \dots x \quad \square$$

Thus if  $k \leq x, y \leq n$  and  $x+y=n$ ,  $f(x)+f(y)$  is minimized for  $x=y=\frac{n}{2}$ . (provided  $k \leq \frac{n}{2}$ ). Thus

$$\begin{aligned}
 p &\geq \frac{\binom{n/2}{k} + \binom{n/2}{k}}{\binom{n}{k}} = \frac{2 \binom{n/2}{k}}{\binom{n}{k}} \\
 &= 2 \frac{(n/2)! / k! (n/2-k)!}{n! / k! (n-k)!} \\
 &= 2 \frac{(n/2-k+1) \dots n/2}{(n-k+1) \dots n} \\
 &= \frac{2}{2^k} \cdot \frac{(n-2(k-1)) \cdot (n-2(k-2)) \dots n}{(n-k-1) \cdot (n-(k-2)) \dots n} \\
 &= \frac{1}{2^{k-1}} \cdot \prod_{i=0}^{k-1} \frac{n-2i}{n-i} \\
 &= \frac{1}{2^{k-1}} \cdot \prod_{i=0}^{k-1} \left(1 - \frac{i}{n-i}\right) \\
 &\geq \frac{1}{2^{k-1}} \prod_{i=0}^{k-1} \exp\left(-\frac{2i}{n-i}\right) \\
 &= \frac{1}{2^{k-1}} \exp\left(-2 \sum_{i=0}^{k-1} \frac{i}{n-i}\right) \\
 &= \frac{1}{2^{k-1}} \exp\left(-\Theta\left(\frac{k^2}{n}\right)\right) \quad (\text{as } n \gg k) \\
 &\geq \Omega\left(\frac{1}{2^k}\right) \quad (\text{maximized for } n = \Theta(k^2))
 \end{aligned}$$

assuming  $n \gg k$   
 using  
 $1-x \geq e^{-2x}$   
 for  $x \in [0, \frac{1}{2}]$

Thus  $m = \frac{n}{p} \ln 2 = \Theta(n 2^k) = \Theta(k^2 2^k)$ .  $\square$